

METHOD FOR NONVOLATILE STORAGE OF AT LEAST ONE OPERATING DATA VALUE  
OF AN ELECTRICAL MOTOR, AND ELECTRICAL MOTOR FOR SAID METHOD

The invention concerns a method for nonvolatile storage of at least one operating data value of an electric motor which comprises a microprocessor or microcontroller, hereinafter called a microprocessor, that controls its commutation, and a nonvolatile memory. It furthermore concerns a motor for carrying out such a method.

It is an object of the invention to make available a method of the kind cited initially, and a corresponding electric motor.

This object is achieved by a method according to claim 1. The result is that the microprocessor responsible for commutation is occupied with further useful operations, namely with the updating of the at least one operating variable; and that said variable needs to be transferred, only at appropriate time intervals, to the nonvolatile memory. Fractions of an hour are, for example, a suitable time interval. The transfer can thus take place relatively infrequently, and therefore does not require a great deal of calculation time. Good accuracy of the data saved in the nonvolatile memory is nevertheless obtained, since electric motors usually run uninterruptedly for a long period during which their operating data change only slightly.

A further preferred approach to achieving the object is evident from a method according to claim 2. Because the at least one operating variable is updated in the time intervals between the commutation operations, the microprocessor is optimally utilized.

Another way of achieving the stated object is evident from the subject matter of claim 9.

Further details and advantageous developments of the invention are evident from the exemplary embodiments - which are in no way to be understood as a limitation of the invention - that are described below and depicted in the drawings, and from the dependent claims. In the drawings:

FIG. 1 is a general overview of an arrangement according to the present invention having an electric motor;

FIG. 2 is an overview circuit diagram of a preferred embodiment having an electronically commutated motor;

FIG. <sup>3A-3C</sup> shows exemplary diagrams of voltages and signals occurring in <sup>10/30/03</sup> ~~2~~ triangular generator 100 of FIG. 2;

FIG. 4 is a schematic diagram of a signal HALL detected by a rotor position sensor 132 and transferred to  $\mu$ C 23 of FIG. 2;

FIG. 5 is an overview diagram of error function 85 and alarm function 86 implemented in FIG. 1;

FIG. 6 shows a main program in the form of a function manager that can preferably be used in the context of a motor according to the present invention;

FIG. 7 depicts a control word having eight bits that serve, in the function manager of FIG. 6, to request the execution of functions or to reset said requests;

FIG. 8 is a table with objects which contain parameters for motor 32 of FIG. 2;

FIG. 9 shows a control word DI\_CTRL that is used by the error function (FIGS. 2 and 23);

FIG. 10 shows a state word DI\_STATE that is used by the error function;

FIG. 11 shows an error code DI\_CODE that is used by the error function;

FIG. 12 shows a reaction word DI\_REAC that is used by the error function;

FIG. 13 is a flow chart of the "Hall Interrupt" function of FIG. 6;  
 FIG. 14 is a flow chart of the "TIMER0 Interrupt" function of FIG. 6;  
 FIG. 15 is a flow chart of the "Operating Data Function" of FIG. 6;  
 FIG. 16 is a flow chart of the "A/D" function of FIG. 6;  
 FIG. 17 is a flow chart of the "Error Detection" function of FIG. 6;  
 FIG. 18 depicts a RAM region of  $\mu$ C 23 of FIG. 2;  
 FIG. 19 is a flow chart of a "CHK\_CALC" function for calculating a  
 check value for memory checking;  
 FIG. 20 is a flow chart of a "RAM\_CHK\_SET" function which serves to  
 calculate a check value RAM\_CHK and save it at a predetermined point in a  
 RAM region;  
 FIG. 21 is a flow chart of the "RAM\_CHK\_TEST" function of FIG. 17,  
 which serves for memory checking;  
 FIG. 22 is a flow chart of a "NEW\_DIST" function which is called in  
 the event of a new error;  
 FIG. 23 is a flow chart of the "Error Function" routine of FIGS. 1  
 and 6; and  
 FIG. 24 is a flow chart of the "COMM" routine of FIG. 6.

In the flow charts hereinafter, Y = Yes and N = No. Identical or  
 identically functioning parts are labeled with the same reference  
 characters and usually are described only once.

GENERAL OVERVIEW (FIG. 1)

FIG. 1 shows an overview of an arrangement according to the present invention having an electric motor 32 that serves, for example, to drive a fan as shown in FIG. 1. Associated with electric motor 32 are a drive function "AF" 150, a temperature sensor 152, an operating data function "BDF" 89, an error function 85, an alarm function "ALARM" 86, a nonvolatile memory 74 (e.g. an EEPROM), a bus 82, a bus interface 80, and an alarm output ALARM\_OUT 88.

Temperature sensor 152 serves to measure temperature T at motor 32 or in its vicinity, in order to control operations as applicable on the basis of said temperature T, for example to define a rotation speed  $n_s$ , dependent on said temperature, of motor 32; or to trigger an alarm signal if said temperature becomes too high; or to save the maximum values of said temperature.

Drive function "AF" 150 ensures that electric motor 32 runs in a desired direction and at the desired rotation speed  $n_s$ . For that purpose, drive function "AF" 150 is connected via a connection 150A to electric motor 32. The drive function can be embodied, for example, as a commutation control system of an electronically commutated motor (ECM).

Drive function 150 is connected to bus 82, which is configured bidirectionally. Bus 82 has an interface 80 to which, for example, a PC 81 can be connected. Via bus 82, drive function "AF" 150, operating data function "BDF" 89, error function 85, and alarm function "ALARM" 86 can be configured, for example, by PC 81 or by another input device. Data from these functions and from EEPROM 74 can, however, also be written onto bus 82 and transferred via it to, for example, the external PC 81 or to another motor: for example, the number of operating hours of motor 32, or a datum concerning errors that have occurred, an extreme temperature, an excessive operating voltage, etc.

Through a connection 154, electric motor 32 delivers operating data to operating data function "BDF". The latter can save the operating data

via a connection 160 in nonvolatile memory 74 ("SAVE DATA"), or write them onto bus 82 ("WRITE BUS") so that said data can be read by PC 81 via bus interface 80.

An error in electric motor 32 is reported, via connection 156, to error function 85. Upon occurrence of an error, the latter can modify rotation speed target value  $n_s$  ("SET  $n_s$ ") via a connection 162 that leads to drive function "AF" 150; write into nonvolatile memory 74, via a connection 166, data concerning the error and instantaneous operating data that are obtainable via a connection 164 from operating data function "BDF" 89 ("SAVE DATA"); output data concerning the error and instantaneous operating data via bus 82 ("WRITE BUS"); or, via a connection 168, trigger an alarm in alarm function 86 ("ALARM").

Alarm function 86 either receives a signal from error function 85 via a connection 168, or receives a signal directly from motor 32 via a connection 158. The alarm function can write data concerning the alarm and instantaneous operating data into nonvolatile memory 74 via a connection 170 ("SAVE DATA"), and/or it writes said data onto bus 82 ("WRITE BUS"), and/or it outputs a signal via a line ALARM\_OUT 88 ("WRITE ALARM\_OUT").

Nonvolatile memory 74 can contain data that were saved by functions 89, 85, 86, and that can also be read out again by them and can be outputted, for example, when polled via bus 82. Nonvolatile memory 74 can moreover contain configuration parameters for functions 150, 89, 85, 86.

Connected to bus 82 is a clock CLK 149 that is backed up by a battery 148 and thus runs continuously. The following models are suitable, for example, for an I<sup>2</sup>C bus: PCF8563, PCF8573, or PCF8583. By means of clock 149, it is possible to save in EEPROM 74 the time of day at which an error occurred, and optionally also the date.

The configuration of a preferred arrangement according to FIG. 1 and of a program used in its context will be described in more detail below.

## MOTOR OVERVIEW (FIG. 2)

FIG. 2 shows an overview of a preferred exemplary embodiment of an electronically commutated motor (ECM) 32 according to the present invention. The latter is controlled by means of a microcontroller ( $\mu$ C) 23, or alternatively a microprocessor.  $\mu$ C 23 has an A/D converter 60, a characteristic function 68, a "RGL" function (controller function) 70, a "CTRL" function (motor current monitor) 71, a "CTRL EEPROM" function 72, a "COMM" function (communication function) 78, "FEHLER" function (error function) 85, "ALARM" function 86, "BDF" function (operating data function) 89, and "AF" function (drive function) 150.

A resistor 52 is connected between a node 56 and ground (GND), and a resistor 54 is present between an operating voltage  $U_b$  and node 56. Node 56 is connected to input 57 of A/D converter 60. This arrangement serves to digitize a value corresponding to operating voltage  $U_b$ .

A Negative Temperature Coefficient (NTC) resistor 62 (in temperature sensor 152) is connected between a node 66 and ground (GND), and a resistor 64 is connected between a regulated voltage  $V_{cc}$  (e.g. +5V) and node 66. Node 66 is connected to input 67 of A/D converter 60. This arrangement serves to digitize a temperature sensed with NTC resistor 62.

EEPROM 74 is connected via a bus 76 to "CTRL EEPROM" function 72. Instead of EEPROM 74 it is possible to use, for example, a flash ROM, a reprogrammable flex ROM line, or another nonvolatile memory. EEPROM 74 can optionally be integrated into  $\mu$ C 23.

Bus interface 80 is connected via bus 82 to communication function COMM 78. The latter is in turn connected via a line 84 to "CTRL-EEPROM" function 72 and to other functional elements of  $\mu$ C 23.

"ALARM" function 86 can output a signal at output ALARM\_OUT 88.

In this embodiment, "RGL" function 70 is, by way of example, connected to a pulse width modulation generator (PWM generator) 100. The PWM generator comprises a control voltage generator 104, a triangular generator 106, and a comparator 102, and its manner of operation is described in more detail with reference to FIG. 3. Through output 107 of PWM generator 100, a signal PWM passes to two logical AND elements 108, 110. The width of pulses 107A of signal PWM is variable.

As a simple example, FIG. 2 depicts an electronically commutated motor 32 having a single phase 128. The principle of a simple motor of this kind is explained, for example, in DE 23 46 380 C. Motor 32 comprises a permanent-magnet rotor 130, a Hall sensor 132, and an output stage 112.

Output stage 112 has four npn transistors 114, 116, 118, 120 connected as an H-bridge, and a low-resistance resistor 124 for current measurement.

A current limiter " $I < I_{max}$ " 125 receives a voltage corresponding to the value of motor current I measured at resistor 124 and, if motor current I is too high, influences control voltage generator 104 so as to reduce said current. Current limiter 125 is also associated with "CTRL" function (motor current monitor) 71 in  $\mu C$  23.

The signal of Hall sensor 132 is delivered to an evaluation circuit 134, which contains a lowpass filter e.g. in the form of a comparator and which generates signal HALL, depicted in FIG. 4, that is delivered to "AF" function 150. The latter controls two outputs OUT1 and OUT2 which control upper transistors 114, 116 and (via AND elements 108, 110) lower transistors 118, 120.

$\mu C$  23 furthermore has a ROM 96, a RAM 97, and a timer 98 which is also referred to as TIMER0. ROM 96 is usually programmed together with the manufacture of  $\mu C$  23. It can also be arranged outside  $\mu C$  23, as can RAM 97 and timer 98, as is known to those skilled in the art.

## MODE OF OPERATION

Phase 128 is energized by means of transistor output stage 112. Outputs OUT1, OUT2 control transistors 114, 116, 118, 120 that are connected as H-bridge 112. If OUT1 is HIGH and OUT2 is LOW, transistors 114 and 118 are conductive and a current flows from operating voltage  $+U_b$  through transistor 114, stator winding 128, transistor 118, and resistor 124 to ground GND. It is assumed in this discussion that signal PWM (line 107) is continuously HIGH, since otherwise AND elements 108, 110 and therefore transistors 118, 120 are blocked.

If OUT1 is LOW and OUT2 is HIGH, a current then flows from  $U_b$  through transistor 116, through stator winding 128 in the opposite direction, and through transistor 120 and resistor 124 to ground GND.

The alternating magnetic flux generated by stator winding 128 exerts a torque on permanent-magnet rotor 130 and drives it. In this exemplary embodiment, rotor 130 is depicted with four poles.

The position of rotor 130 is sensed via Hall sensor 132. The latter's signal is filtered through a lowpass filter in circuit 134 and processed to yield a square-wave digital signal HALL (FIG. 4), which is delivered to "AF" function 150.

"AF" function 150 controls outputs OUT1, OUT2 on the basis of signal HALL. In this example, the commutation of motor 32 is accomplished electronically and is explained below with reference to FIG. 13. "AF" function 150 moreover ensures correct commutation for reliable operation of motor 32, for example in the event of an overload of transistor output stage 112. Commutation can also be implemented in such a way that transistors 114 through 120 are commutated earlier as the rotation speed increases, somewhat analogously to ignition advance in a gasoline engine.



Of course the invention is similarly suitable for any type of motor, e.g. for three-phase ECMs and others. This is therefore only a simple exemplary embodiment in order to facilitate understanding of the invention.

In this exemplary embodiment, rotation speed control is achieved by controlling pulse duty factor PWM\_TV of signal PWM at output 107 of controller 100, i.e. by making pulses 107A longer or shorter (see FIG. 3C for a definition of pulse duty factor PWM\_TV). The greater this pulse duty factor, the longer pulses 107A become, and the longer the output of AND element 108 or 110 (controlled at the time by OUT1 or OUT2) is switched to HIGH. Stator winding 128 is therefore energized for a longer period, and motor 32 is more strongly driven. If OUT1 is HIGH and OUT2 is LOW, for example, upper transistor 114 is made conductive, and lower transistor 118 is switched on and off by AND element 108 in accordance with signal PWM.

In this exemplary embodiment, "RGL" function 70 controls rotation speed  $n$  of motor 32 via PWM generator 100. For that purpose, "RGL" function 70 has available to it rotation speed  $n$  of rotor 130, which is calculated using signal HALL (see description with reference to FIG. 4), and rotation speed target value  $n_s$ , which in this exemplary embodiment is determined by characteristic function 68. Rotation speeds  $n$  and  $n_s$  can be present, for example, in the form of Hall times  $t_H$  (FIG. 4), e.g. in units of microseconds or seconds; or as rotation speeds, e.g. in units of revolutions per minute.

In this example, characteristic function 68 associates a rotation speed target value  $n_s(T)$  with each temperature  $T$  sensed by sensor 152 of FIG. 1. Temperature  $T$  is sensed by means of NTC resistor 62 which, together with resistor 64, constitutes a voltage divider between  $V_{cc}$  and ground (cf. FIG. 2). The potential at node 66, which constitutes an indication of the temperature of resistor 62, is digitized by A/D converter 60 located in  $\mu C$  23, and delivered to characteristic function 68.

From temperature T, characteristic function 68 determines rotation speed target value  $n_s(T)$  of motor 32. For that purpose the value  $n_s(T)$  pertinent to temperature T is loaded, for example via a "CTRL EEPROM" function 72, from a temperature/rotation speed target value table in EEPROM 74.

The "COMM" function manages bus interface 80 over which data can be transferred from outside into  $\mu C$  23, and over which, conversely, data can be transmitted to the outside from  $\mu C$  23. For example, data that arrive in  $\mu C$  23 via bus interface 80 by means of "COMM" function 78 can be written, via connection 84 and by means of "CTRL EEPROM" function 72, into EEPROM 74.

#### PWM GENERATOR (FIG. 3)

FIG. 3A shows a triangular signal u106 of triangular generator 106, and a control output u104 that is generated by control voltage generator 104. FIG. 3B shows pulses 107A resulting from FIG. 3A, and FIG. 3C shows the calculation of pulse duty factor PWM\_TV of pulses 107A.

Triangular signal u106 from triangular generator 106 is depicted in idealized fashion. In reality it is not perfectly triangular in shape, although this makes no difference in terms of the manner of operation of PWM generator 100 of FIG. 2. Triangular signal u106 has an offset 139 from voltage 0V. Control output u104 therefore causes a pulse duty factor PWM\_TV > 0 only when it is greater than offset 139.

Pulse duty factor PWM\_TV of signal PWM is the ratio between the duration  $t_{ON}$  that signal PWM is HIGH during one period of triangular signal u106, and one entire period T of triangular signal u106 (cf. FIG. 3B). The equation is:

$$PWM\_TV = t_{ON} / T \quad (1)$$

Pulse duty factor PWM\_TV can be between 0 and 100%. For example, if the motor rotation speed is too high, control output u104 is then decreased and pulse duty factor PWM\_TV is thus made smaller, as depicted in FIG. 3A. This entire procedure is referred to as pulse width modulation (PWM).

A motor 32 according to the present invention can, of course, also be operated without pulse width modulation, e.g. without control or with a different kind of control. This serves only as an example to facilitate comprehension.

SIGNAL HALL (FIG. 4)

FIG. 4 shows signal HALL, which corresponds to the position of rotor 130 sensed by Hall sensor 132 (FIG. 1) and is delivered to  $\mu C$  23 via circuit 134 (FIG. 2).

As an example, rotor 130 can have a rotation speed  $n = 6000$  rpm, corresponding to 100 revolutions per second. One mechanical rotation of rotor 130 then lasts 10 ms. Rotor 130 is depicted with four poles in this exemplary embodiment, so that four Hall changes - two from HIGH to LOW and two from LOW to HIGH - take place in one mechanical revolution ( $360^\circ$  mech.). One electrical revolution ( $360^\circ$  el.), on the other hand, has already taken place after only two Hall changes. In a four-pole motor, therefore, two electrical revolutions take place for one mechanical revolution.

Rotation speed  $n$  is calculated from Hall time  $t_H$  (FIG. 4) between two Hall changes, using

$$t_H = T / P \quad (2)$$

In addition,  $T = (60 \text{ seconds}) / n \quad (3)$

Combining (2) and (3) yields

$$t_H = ((60 \text{ seconds}) / n) / P \quad (4)$$

where

$T$  = duration (in seconds) of one mechanical revolution of rotor 130;

$P$  = number of poles of the rotor (here  $P = 4$ ); and

$n$  = rotation speed in rpm.

If  $n = 6000$  rpm and  $P = 4$ , equation (4) yields

$$t_H = 60 \text{ seconds} / 6000 / 4 = 2.5 \text{ ms.}$$

At a rotation speed of 6000 rpm, time offset  $t_H$  between two changes of signal HALL is therefore 2.5 ms, as depicted by way of example in FIG. 4.

#### OVERVIEW OF ERROR AND ALARM FUNCTIONS (FIG. 5)

FIG. 5 shows an overview of the interaction, in terms of program engineering, between error function 85 and "ALARM" function 86 for a motor 32 as shown in FIG. 1.

"Sensor interruption" function 91, "bus error check" function 92, "temperature check" function 94, and "rotation speed check" function 95 are depicted in the top row. The number 93 represents any further checks that are not depicted.

Considered in terms of program engineering, the checks are located at the point at which the respective measurement takes place. When A/D converter 60 is polled, for example, a check is made on the basis of the value in "sensor interruption check" 91 as to whether a "sensor interruption" of NTC resistor 62 (FIG. 1) is present (cf. FIG. 16, S226 below). This means that in FIG. 2, the connection to NTC resistor 62 is interrupted at point 62a and/or 62b, i.e. the line has been broken. In this case "sensor interruption check" 91 reports an error, i.e. it generates an error signal. This is subsequently noted by "error function" 85 (FIG. 23), which then decides what else will happen. For example, it can set rotation speed target value  $n_s$  to a maximum value and request an alarm from alarm function 86. This is described in detail below with reference to FIG. 23, and happens on the basis of parameters that are saved in EEPROM 74 and can be modified.

"Temperature check" 94 and "rotation speed check" 95 occupy a special place. Temperature and rotation speed are so important for the functionality of a fan that an excessive deviation in rotation speed, or the fact that a predetermined temperature has been exceeded, is forwarded directly to "alarm function" 86.

Other errors, however, for example a sensor interruption in NTC resistor 62, do not rule out satisfactory operation of the fan

and can therefore be processed by error function 85 (FIG. 23). Error function 85 can be parameterized as desired by the customer, as will be described later.

#### FUNCTION MANAGER (FIGS. 6 and 7)

FIG. 6 shows a flow chart with one possible embodiment of the main program executing in  $\mu$ C 23, in the form of a so-called function manager 601.

The task of the main program is to react to events, e.g. to a change in signal HALL; also to make resources, in particular calculation time, available to each function as necessary; and to observe priorities in assigning resources.

After motor 32 is switched on, an internal reset is triggered in  $\mu$ C 23, and initialization of  $\mu$ C 23 takes place in S600. In this context, data are loaded from EEPROM 74 into RAM 97 of  $\mu$ C 23 so that they are quickly available for program execution. A memory test is also accomplished, as described below.

After initialization, execution jumps into function manager 601, which begins in S602. Those functions that are time-critical and must be executed at each pass are executed first. These include functions "COMM" in S604 (cf. FIG. 24), "A/D" in S606 (cf. FIG. 16), "I\_max" in S608, and "RGL" in S610.

In "COMM" function (S604), communication via bus 82 (FIG. 1) is monitored. At a baud rate of, for example, 2 K, bus 82 must be checked every 250 microseconds.

In S606, A/D converter 60 (FIG. 2) is polled. It digitizes the potentials at inputs 57, 67. Further A/D converters for digitizing further potentials can be present.

In S608, an "I\_max" motor current limiting routine that may be present is executed.

The "RGL" function for controlling rotation speed n is called in S610.

FIG. 7 shows an example of a function register 605 in which one bit is reserved for each of the functions in S622, S626, S630, S634, and S638 (FIG. 6).

In this example, function register 605 is one byte long, and the following bits, beginning at the least significant bit (LSB), are defined for the requestable functions explained below:

- FCT\_KL        for the characteristic function
- FCT\_n        for the rotation speed calculation function
- FCT\_AL\_n     for the alarm rotation speed check
- FCT\_DIST     for error detection
- FCT\_BDF      for the operating data function.

The remaining bits are reserved for additional requestable functions that may be inserted into the function manager as necessary.

If, in FIGS. 6 and 7, a specific requestable function is to be requested by another function or by an interrupt routine, the bit of the function to be requested is set in function register 605 to 1, for example  $FCT\_AL\_n := 1$ . If function manager 601 (FIG. 6) then finds, at the pass following this request, no other requestable function with a higher priority, the aforesaid function (i.e. the alarm rotation speed check) is therefore called in S630.

Once a requested function has been executed, it sets its bit in function register (FIG. 7) back to 0, (e.g.  $FCT\_AL\_n := 0$ ) at the end of S630.

Once the requestable function has been performed, execution jumps back to S602 at the beginning ("FCT\_MAN") of function manager 601.

In FIG. 6, after S610 the program begins with the most important requestable function and checks in a predetermined sequence as to whether its request bit is set. If so, the requested function is then performed.

The higher up such a function is located in function manager 601, the higher its priority.

S620 checks whether request bit FCT\_KL is set. If it is set, the characteristic function is called in S622.

If FCT\_n is set in S624, the rotation speed calculation function is called in S626.

If FCT\_AL\_n is set in S628, the alarm rotation speed check is called in S630.

If FCT\_DIST is set in S632, error detection (described with reference to FIG. 17) is called in S634.

If FCT\_BDF is set in S636, the operating data function (described with reference to FIG. 15) is called in S638.

If none of the request bits of function register 605 was set, an error function is executed in S640 and an alarm function in S642, and execution branches back to S602. See FIG. 23 regarding error function S640; it can also be referred to as an error monitoring routine, since it monitors whether any of the other routines has reported an error, and then implements a reaction to that error.

FIG. 6 also symbolically shows a Hall interrupt 611 (FIG. 13), which has the highest priority L1 (level 1). It interrupts all the processes of function manager 601, as symbolized by an arrow 613, in order to achieve precise commutation of motor 32. A Hall interrupt 611 is generated each time signal HALL changes in FIG. 4, and it causes an incrementing of commutation counter CNT\_COM as described with reference to FIG. 13. Commutation of motor 32, i.e. the generation of signals OUT1 and OUT2, is also controlled directly or indirectly by the Hall interrupts 611 in order to make the motor run smoothly (cf. also FIG. 13).

Depicted below Hall interrupt 611 at 615 is a TIMERØ interrupt of

timer TIMER0 98 (FIG. 2). It has a lower priority L2 and interrupts all processes below it, as indicated by arrow 617. It is described with reference to FIG. 14.

If Hall interrupt 611 and TIMER0 interrupt 615 were requested simultaneously, they would be executed in the order of their priority.

The subsequent functions have increasingly lower priorities, from L3 for the "COMM" function in S604 to L13 for the alarm function in S642.

In this fashion, it is possible to classify the various "needs" of motor 32 into a predetermined hierarchy, and to use the resources of  $\mu$ C 23 optimally for the operation of motor 32. Error function S640 and alarm S642 are thus executed only when  $\mu$ C 23 presently has free calculation time.

OBJECT TABLE (FIG. 8)

FIG. 8 shows a table 111 with objects that contain configuration parameters for motor 32. The individual objects have an index, a memory type, access rights, and a name.

Object table 111 is saved in a nonvolatile memory, in this exemplary embodiment in EEPROM 74 (FIG. 1). After each reset of  $\mu$ C 23, upon initialization in S600 (FIG. 6) object table 111 is transferred by "CTRL EEPROM" function 72 out of EEPROM 74 into RAM 97 of  $\mu$ C 23, and is thereupon available to the program (FIG. 6) executing in  $\mu$ C 23.

The index in table 111 is indicated hexadecimally, a "0x" before a number indicating hexadecimal notation. The memory type is either "unsigned8" (one byte with no sign bit), "unsigned16" (two bytes with no sign bit), or "unsigned24" (three bytes with no sign bit). The access rights are R (read) and W (write). The objects can be read out and modified. The name of the object makes utilization easier. The names denote:



DIST_CTRL	Control word for the error function
DIST_STATE	State word for the error function
DIST_CODE	Error code for the error function
DIST_REAC	Reaction word for the error function
n_DIST	Error rotation speed
t_COMM_TO	Maximum time-out time for the communication function
OD_TMAX	Temperature for the operating data function
OD_UBMAX	Operating voltage for the operating data function
OD_OHO	Operating hours (e.g. in units of 10 minutes) for the operating data function
OD_COMMUT	Total number of commutations (e.g. in units of 10,000) for the operating data function

Because of the open structure of object table 111, it is easily possible to insert new objects using a standardized procedure and to expand the table as desired. Any modification of object table 111 and thus of the configuration is preferably accomplished via bus 82, "COMM" function 78 (FIG. 2), and "CTRL EEPROM" function 72. Configuration can be performed to the customer's specification before delivery, or customers can be given the capability to make modifications themselves.

#### EXPLANATION OF THE OBJECTS IN OBJECT TABLE 111

DIST\_CTRL is the control word for error monitoring routine 85 that is labeled S640 in FIGS. 6 and 23. Its structure is evident from FIG. 9 and the accompanying description. Depending on its content, this word causes an error either to be saved in EEPROM 74 or not to be saved. When DIST\_CTRL is loaded into RAM 97, it is labeled DI\_CTRL.

DIST\_STATE is a state word for the error function. When DIST\_STATE is loaded into RAM 97, it is labeled DI\_STATE. The structure of DI\_STATE is evident from FIG. 10 and the accompanying description. It indicates in its bit 7 whether an error is present, and its bits 0 through 2 approximately define the type of error that has occurred, e.g. an error in commutation or an error in sensor 152.

DIST\_CODE contains error codes for error function 85 which specify the error more precisely. When DIST\_CODE is loaded in RAM 97, it is labeled DI\_CODE. Its structure is evident from FIG. 11 and the accompanying description. There can be, for example, four separate error codes concerning the type of transfer error for error class DS\_COMM of state word DI\_STATE.

DIST\_REAC is a reaction word for error function 85, and indicates how the motor is to react to an error, e.g. by stopping or braking, or with maximum rotation speed. When DIST\_REAC is loaded into RAM 97, it is labeled DI\_REAC. Its structure is evident from FIG. 12 and the accompanying description. FIG. 23 shows how it is evaluated.

n\_DIST is the error rotation speed. This is a fixed rotation speed at which motor 32 is to run in the event of an error (cf. FIG. 23, S378 and S380).

t\_COMM\_TO is the maximum time-out time for the communication function. It determines the transfer rate on bus 82.

OD\_TMAX is the extreme upper value of temperature T measured by sensor 152. This value operates in exactly the same way as a "maximum thermometer," but digitally. Every 10 minutes, a check is made as to whether the present temperature is higher than the saved value OD\_TMAX; if so, the new, higher value is saved in EEPROM 74 as OD\_TMAX. This can be important for the analysis of errors.

OD\_UBMAX is the extreme upper value of operating voltage UB of motor 32. Every 10 minutes, a check is made as to whether the present operating voltage UB is higher than the saved value OD\_UBMAX; if so, the new, higher value is saved in EEPROM 74 as OD\_UBMAX. This can be important for the analysis of errors.

OD\_OHO is the total number of operating hours of motor 32, measured in units of 10 minutes; OD\_OHO = 6,000 thus means 1,000 operating hours.

At startup, OD\_OHO is loaded into RAM 97 and continuously updated therein by means of the routine shown in FIG. 14. Every 10 minutes the updated value is written into EEPROM 74 together with other values, as explained with reference to FIG. 15.

OD\_COMMUT is the total number of Hall interrupts 611 (FIG. 4). Since each Hall interrupt causes a commutation (cf. FIG. 13), this is an indication of the total number of revolutions of rotor 130.

The number of commutations is saved in units of 10,000, so a value of 200,000 indicates that 2 billion commutations have taken place. With a four-pole rotor 130 this corresponds to 500 million revolutions, and at a constant 3,000 rpm this would correspond to an operating time of approximately 2,800 hours. This figure provides information about the expected remaining service life of the bearings of motor 32.

At startup, OD\_COMMUT is loaded into RAM 97 and is continuously updated therein by means of the routine shown in FIG. 13. Every 10 minutes the updated value is written into EEPROM 74 together with other values, as explained with reference to FIG. 15.

Upon occurrence of an error, the present values (in RAM 97) of OD\_TMAX, OD\_UBMAX, OD\_OHO, and OD\_COMMUT are saved in a FIFO of EEPROM 74 (cf. FIG. 22, S346, push\_OD\_DATA).

Object table 111 (FIG. 8) thus contains a kind of "curriculum vitae" of motor 32; and, by modifying the first four objects, it is possible to define whether and how motor 32 reacts to an error, what is to be saved and how, where an alarm will be outputted, etc. In other words, corresponding parameters can be defined and entered into the motor.

#### CONTROL WORDS AND STATE WORDS

FIG. 9 shows a control word DI\_CTRL, which has the same structure as object DIST\_CTRL from object table 111 (FIG. 8) and which serves for data exchange between the operating system (FIG. 6) and the individual error

functions. It is located in RAM 97, into which the value of DIST\_CTRL (FIG. 8) is loaded in "INIT" S600 at the start of the main program after a reset of  $\mu$ C 23. The bits of DI\_CTRL are serially numbered 0 through 7.

Bit 0 is named DC\_LATCH. If DC\_LATCH = NO\_LATCH, i.e. LOW (0), an error is not saved, and the error state is reset after the error disappears. If DC\_LATCH = LATCH, however, i.e. HIGH (1), an error state is not cleared until a request to do so comes via bus 82.

Bits 1 through 6 are not used here, and are reserved (RES) for future uses and enhancements.

Bit 7 is named DC\_CLEAR, and is used to clear an error state. To do so, the value of DC\_CLEAR, which is 0 in the base state, must be set to 1 and then back to 0, as indicated in FIG. 9.

FIG. 10 shows a state word DI\_STATE, which corresponds to object DIST\_STATE (FIG. 8) and is initialized in "INIT" S600 in the same way as DI\_CTRL.

Bits 0 through 2 are named DS\_CLASS, and DS\_CLASS can assume the (decimal) values 0 through 7. DS\_CLASS contains the error class. The error classes are defined as follows:

- DS\_μC (0)      Error in  $\mu$ C 23
- DS\_COMM (1)      Communication error
- DS\_SENS (2)      Error in the sensor or sensors, e.g. NTC resistor
- 62
- DS\_HW (3)      Error elsewhere in the hardware.

The remaining values 4 through 7 of DS\_CLASS are not used (RES) in this exemplary embodiment.

Bits 3 through 6 are not used (RES).

Bit 7 is named DS\_ACTIVE. If DS\_ACTIVE = NO\_DIST (0), no error is present and the content of DS\_CLASS is irrelevant. If DS\_ACTIVE = DIST (1), an error is present and DS\_CLASS contains the error class.

FIG. 11 shows an error code DI\_CODE which specifies the error more precisely. After each reset of  $\mu$ C 23, in S600 (FIG. 6) the value of object DIST\_CODE (FIG. 8) is written to DI\_CODE. DI\_CODE is 16 bits long, and can therefore represent values from 0 to 65535. One thousand values are provided for each of the individual error classes (FIG. 10).

Values (VAL) 0 through 999 are reserved for class DS\_ $\mu$ C, i.e. for errors of  $\mu$ C 23. The error code definitions are as follows:

- DN\_WDT                      Error in a watchdog timer
- DN\_CHKS\_ROM                Checksum error in ROM 96
- DN\_CHKS\_RAM                Checksum error in static portion 142 of RAM 97
- DN\_CHKS\_EEPROM            Checksum error in EEPROM 74
- DN\_TEST\_RAM                Error in internal RAM test of  $\mu$ C 23.

Values 1000 through 1999 are reserved for class DS\_COMM. The error code definitions are as follows:

- DN\_TIMEOUT\_TRANSFER        Time-out error during a transfer
- DN\_TIMEOUT\_BUS            Time-out error during access to bus 82
- DN\_PROT\_ERR                Invalid transfer protocol (e.g. 9 data bits)
- DN\_INVAL\_DATA              Invalid data

Values 2000 through 2999 are reserved for class DS\_SENSOR. The error code definitions are as follows:

- DN\_SENSOR\_INTERRUPT        Sensor interruption (interruption at point 62a or 62b of FIG. 2)
- DN\_SENSOR\_SHORT            Sensor short circuit (between points 62a and 62b)

Values 3000 through 3999 are reserved for class DS\_HW (HW = hardware). The error code definitions are as follows:

- DN\_DRIVER\_FAULT            Error in output stage 112 (FIG. 2).

The remaining values are not defined here (RES).

FIG. 12 shows a reaction word DI\_REAC which indicates the reaction that should occur in response to an error. DI\_REAC corresponds to object DIST\_REAC of FIG. 8 and, like DI\_CTRL, it is written out from object table 111 in S600 (FIG. 6), along with the value of object DIST\_REAC, after a reset of  $\mu$ C 23.

Bits 0 through 2 are named DR\_REAC, and DR\_REAC can assume a (decimal) value (VAL) of 0 through 7. DR\_REAC contains the reaction by the error function in response to an error. The reactions are defined as follows:

- DR\_OFF            No reaction
- DR\_n\_max        Maximum rotation speed
- DR\_n\_min        Minimum rotation speed
- DR\_n\_0          Zero rotation speed
- DR\_n\_DIST       Specific error rotation speed
- DR\_BRAKE        Zero rotation speed and active braking of motor 32.

Bit 3 is named DR\_AL. If DR\_AL = DR\_AL\_OFF, no alarm is requested from alarm function 86 in the event of an error. If, on the other hand, DR\_AL = DR\_AL\_ON, then an alarm is requested from alarm function 86 in the event of an error. Bits 4 through 7 are not used here (RES).

#### HALL INTERRUPT AND TIMER0 INTERRUPT

FIG. 13 shows the portions essential in this context of a "Hall interrupt" routine S147 that is called upon occurrence of a Hall interrupt (611 in FIG. 6). A Hall interrupt 611 is triggered at each change in signal HALL (FIG. 4) from HIGH to LOW or from LOW to HIGH, i.e. at times  $t = 0, 2.5, 5, 7.5,$  and  $10$  ms in the example of FIG. 4.

S151 is a general designation of steps that pertain to the calculation of HALL time  $t_H$  (FIG. 4), e.g. stopping a corresponding timer, etc.

In steps S153, S155, and S157, the edge of signal HALL at which the next Hall interrupt is to be triggered in  $\mu$ C 23 is set. For that purpose, S153 checks whether HALL = 1. If Yes, in S155 the edge at which the next

Hall interrupt is to be triggered is set to a trailing edge (HIGH -> LOW). If No, then in S157 the resolution is set to a leading edge (LOW -> HIGH).

In S159, OUT1 and OUT2 are set to zero, i.e. motor 32 is made currentless. The purpose of this is to interrupt H-bridge 112 briefly, so that a short circuit cannot occur in it during a commutation.

A variety of steps can be performed in S159A, e.g. restarting of a counter (not depicted) for the measurement of  $t_H$ . These program steps should last, for example, 50 microseconds.

In S161 through S165, commutation is performed. If HALL = 1 in S161, then in S163 OUT1 is set to 1 while OUT2 remains at 0 (cf. S159). If HALL = 0 in S161, then in S165 OUT2 is set to 1 while OUT1 remains at 0 (cf. S159).

The signal OUT1=1 causes transistors 114 and 118 to be switched on, as already described; and the signal OUT2=1 causes transistors 116 and 120 to be switched on.

Steps S167 through S171 represent a counter with which a counter OD\_COM, which is loaded after each reset of  $\mu C$  23 in S600 (FIG. 6) with the value of object OD\_COMMUT from object table 111, is incremented by 1 e.g. every 10,000 commutations.

To achieve this, in S167 a counter CNT\_COM is incremented by 1 at each Hall interrupt. S169 checks whether CNT\_COM > 9999. If Yes, then in S171 CNT\_COM is set to 1, and counter OD\_COM is incremented by 1. If CNT\_COM is not greater than 9999 in S169, execution then branches directly to the end at S172.

The number of commutations is required in FIG. 6 for the operating data function S638 and error function S640.

FIG. 14 shows a portion of "TIMER0 interrupt" routine S173, which is

labeled 615 in FIG. 6. Every 256 microseconds, for example, timer TIMER0 triggers a TIMER0 interrupt. TIMER0 can therefore be used for time measurements.

Step S174 represents any other applications of timer 98 that are not depicted here.

Steps S175 through S180 represent a subtimer that executes a step S180, for example, every 10 minutes. A timer CNT\_TI, which is incremented by 1 at each TIMER0 interrupt in S176, is used for this purpose. S178 checks whether CNT\_TI > 2,399,999. If Yes, execution branches to S180; otherwise it branches to the end S182.

If CNT\_TI has reached a value of 2,400,000, this means that 256 microseconds have elapsed 2,400,000 times. This corresponds to exactly 10 minutes. CNT\_TI is then reset back to 1 in S180. An "operating hour counter" OD\_OH in RAM 97, to which the value of object OD\_OHO is written from EEPROM 74 after each reset of  $\mu$ C 23, is incremented by 1; and bit FCT\_BDF of function register 605 of FIG. 6 is set, so that "operating data function" S638 is called by function manager 601 (cf. FIG. 6) in order to load specific operating data values into EEPROM 74.

"Operating hour counter" OD\_OH in RAM 97 therefore contains, in this case, the total operating time of the fan in units of 10 minutes, as does operating hour counter OD\_OHO in EEPROM 74.

OPERATING DATA FUNCTION BDF (FIG. 15)

FIG. 15 shows "operating data function" S638 (FIG. 6), which is called when function manager 601 (FIG. 6) reaches step S636 and bit FCT\_BDF of function register 605 (FIG. 7) has been set, for example by the subtimer in the "TIMER0 interrupt" routine S170 (FIG. 14). In the present exemplary embodiment, this happens every 10 minutes.

"Operating data function" S638 serves to save important operating data values (e.g. maximum fan temperature or operating hours) in EEPROM 74, for example in order to obtain a criterion for replacement of the fan.



These data can then be read out of EEPROM 74 and furnish a kind of "health bulletin" about motor 32.

"Disease bulletins" about motor 32 can additionally be saved in a FIFO of EEPROM 74 if such "diseases" occur. These disease bulletins can also be read out of the EEPROM later, and furnish a kind of log of the error that occurred and, optionally, its cause, e.g. excessive temperature, overvoltage, or end of service life.

In this embodiment of operating data function BDF, S190 checks whether the present operating voltage U<sub>B</sub> which was digitized by A/D converter 60 is greater than the previous highest operating voltage OD\_UBM. If Yes, in S192 U<sub>B</sub> is assigned to the value OD\_UBM, and object OD\_UBMAX, which in this exemplary embodiment is located at point pOD\_UBMAX in EEPROM 74, is set to the new value OD\_UBM using the instruction write\_EE. The point at which an object is located in EEPROM 74 (cf. FIG. 8) is designated by a "p" prefix. For example, pOD\_UBMAX is the point in EEPROM 74 at which object OD\_UBMAX is located (cf. FIG. 8). The maximum operating voltage UBMAX is important because too high an operating voltage accelerates wear on the electronic components. It would similarly be possible to retain a minimum operating voltage (undervoltage).

S194 checks whether the present temperature T, which is measured with sensor 152, is greater than the previous maximum temperature OD\_TM. If Yes, then in S196 the new maximum temperature is assigned to the previous maximum temperature OD\_TM, and object OD\_TMAX in object table 111 (FIG. 8) is, by analogy with S192, overwritten with the new maximum value OD\_TM using the instruction write\_EE. (The parameter pOD\_MAX once again indicates the point in EEPROM 74 at which object OD\_TMAX is stored.)

In S198, the present operating hours located in OD\_OH are written, using the instruction write\_EE, into object OD\_OHO (FIG. 8) of EEPROM 74, so that they are retained even when fan 32 is switched off.

The measurement of operating hours OD\_OH is explained with reference to FIG. 14.

In S200 the present number of commutations, which is stored in OD\_COM, is written into object OD\_COMMUT (FIG. 8) using the instruction write\_EE. Measurement of the number of commutations is explained with reference to FIG. 13.

In S202, request bit FCT\_BDF is reset to zero since "operating data function" S638 is completely executed, and function S638 ends at S204.

#### ERROR DETECTION

FIG. 16 shows a flow chart with a portion of "A/D" function S606 (FIG. 6).

In S220, the potential at input 57 of A/D converter 60 (FIG. 2) is read in using the instruction AD(AD\_UB), and is saved in U\_B. The value U\_B corresponds to the present operating voltage UB, e.g. 40 V.

In S222, the potential at input 67 of A/D converter 60 (FIG. 6) is read in using the instruction AD(AD\_T), and is saved in T. The value T corresponds to a present temperature at NTC resistor 62, e.g. 84°C.

Any further steps, for example a request for "characteristic function" S622, are performed in S224.

S226 checks whether a sensor interruption is present, i.e. whether the connection to NTC resistor 62 is interrupted at point 62a or 62b. This is the case if the value for T is less than a sensor interruption value T\_SI. If so, at S228 temperature value T is set to T\_SI, and the program prepares to call a NEW\_DIST function which is explained with reference to FIG. 22. For this purpose, error class DS\_μC and error code DN\_SENSOR\_INTERRUPT for the sensor short-circuit error are saved into variables TEMP\_CLASS and TEMP\_CODE, respectively, and NEW\_DIST (FIG. 22) is called in S230.

S232 checks analogously for a sensor short circuit, i.e. a short circuit between points 62a and 62b in FIG. 2. This is done by determining whether value T is greater than a sensor short-circuit value T\_SS. If Yes, then in S234 value T is set to sensor short-circuit value T\_SS, variables TEMP\_CLASS and TEMP\_CODE are set to the values for a sensor short circuit, and NEW\_DIST (FIG. 22) is called in S236.

Further steps follow, if applicable, in S238. The A/D routine ends at S240.

Instead of or in addition to operating voltage U\_B, a different voltage, e.g. a 12V auxiliary voltage being used, could also be measured in order to save its extreme value.

FIG. 17 is a flow chart for "error detection" function S634 (FIG. 6). This is a requestable function that must be requested using request bit FCT\_DIST=1. It is requested upon initialization (S600 in FIG. 6) after the motor is switched on, and thereafter approximately every 100 ms by a timing member, e.g. a counter controlled by TIMER0. The effect of the error detection function is therefore that the memories are checked every 100 ms.

In S272 a "RAM\_CHK\_TEST" function is performed, in S274 a "ROM\_CHK\_TEST" function, and in S276 a "EEPROM\_CHK\_TEST" function. These functions check whether an error has occurred in RAM 97, ROM 96, or EEPROM 74. It is possible, for example, for a bit in RAM 97 to "flip"; this can lead to errors in the program of  $\mu$ C 23 and thus to unreliable operation of fan 32.

In S277, request bit FCT\_DIST is reset to zero since "error detection" function S634 has been completely executed.

The memory tests will be explained with reference to the subsequent Figures, using the example of the "RAM\_CHK\_TEST" function for testing RAM 97. The tests for EEPROM 74 and ROM 96 proceed in entirely analogous fashion.

FIG. 18 shows a region 140 of RAM 97, which in this case is divided into a static region "STATIC" 142 and a non-static (dynamic) region "NON\_STATIC".

For illustration, memory words W1 through W11 are sequentially numbered. Region "STATIC" 142 comprises memory words W1 through W7. Memory word W8 contains variable RAM\_CHK 146 which serves to check region "STATIC" 142. Region "STATIC" 142 contains, for example, functions and constants. Region "NON\_STATIC" 144 comprises memory words W9 through W11 and contains, for example, variables.

FIG. 19 shows a "CHK\_CALC" function S290 which serves to calculate a check value CHK. For this, in S292 CHK is set to zero and a loop counter N is also set to zero.

A loop begins in S294. N is incremented by 1 each time, and an XOR operation on CHK and RAM(N) is performed. RAM(N) is the memory word in RAM region 140 (FIG. 18) at point WN. S296 checks whether N < 7. If No, this means an XOR of CHK with memory words W1 through W7, and thus with all the memory words in region "STATIC" 142, has been made, and the routine goes to step S298 (END).

FIG. 20 shows a flow chart for "RAM\_CHK\_SET" function S300 which is used to set RAM\_CHK 146 (FIG. 18).

The "CHK\_CALC" function (FIG. 19), which calculates check value CHK, is called in S302. In S304 this value is saved in RAM\_CHK 146 (FIG. 18), i.e. in word W8 of RAM region 140. The routine ends at S306.

The "RAM\_CHK\_SET" function is called, for example, after each reset of  $\mu$ C 23 in S600 (FIG. 6). If a memory word in region "STATIC" 142 (FIG. 18) is deliberately modified during the runtime, "RAM\_CHK\_SET" S300 must be

called again. In ROM 96, a corresponding ROM\_CHK check word is calculated and entered before ROM 96 is "burned."

FIG. 21 is a flow chart of the "RAM\_CHK\_TEST" function (S272 in FIG. 17). In S312, the "CHK\_CALC" function (S290 in FIG. 19) is called again, and check value CHK is calculated.

In S314, CHK is compared to the saved value RAM\_CHK 146 (FIG. 18). If the two values are not equal, an error has occurred in region "STATIC" 142, and execution branches to S316. In S316, variable TEMP\_CLASS is loaded with DS\_μC and variable TEMP\_CODE is loaded with DN\_CHKS\_RAM, and in S318 the "NEW\_DIST" function (FIG. 22) is called.

If the values RAM\_CHK and CHK are identical in S314, the routine branches directly to S320 (END).

The memory checks of ROM 96 and EEPROM 74 proceed analogously. Instead of the XOR procedure it is also possible to use, for example, a checksum method or another check method, e.g. CRC (cyclic redundancy check).

FIG. 22 shows a flow chart of "NEW\_DIST" function S340 which is called each time an error is detected.

S342 checks whether state word bit DS\_ACTIVE = 1 (cf. FIG. 10). If Yes, then an error is already present; execution branches to the end S348 and the new error is ignored. The reason for this is that one error can lead to consequential errors; the first, oldest error is therefore the most important for analysis.

If S342 finds that a new error is present (DS\_ACTIVE = 0), execution then branches to S344. In S344 DS\_ACTIVE (FIG. 10) is set to 1, and error class DS\_CLASS and error code DI\_CODE are set to the respective values

TEMP\_CLASS and TEMP\_CODE set by the calling function. FIG. 21 shows, for example, that TEMP\_CLASS has been set to DS\_μC.

In S346, error class DS\_CLASS, error code DI\_CODE, and the present operating data OD\_DATA are saved, by means of an instruction push\_FIFO, in a FIFO in EEPROM 74. OD\_DATA can contain, for example, the present temperature, present operating hours, present number of commutations, error class, error code, and - if a real-time clock (FIG. 1) is present - the present time of day and date.

Processor 23 then sets a state signal that is continuously checked by a PC 81 connected thereto and causes the latter to call the saved data regarding errors. If motor 32 is stationary as a result of the errors, e.g. because its rotor 130 is jammed, PC 81 can switch on a reserve motor (not depicted), or an alarm is triggered. If motor 32 is a fan, PC 81 can switch another fan to a higher rotation speed so that cooling continues to be guaranteed.

#### ERROR FUNCTION (ERROR MONITORING ROUTINE)

FIG. 23 is a flow chart of error function S640 (FIG. 6), which can also be referred to as the error monitoring routine.

S362 checks or monitors, on the basis of DS\_ACTIVE (FIG. 10), whether an error is present. If DS\_ACTIVE = 0, no error is present; and execution branches to the end (S390).

If DS\_ACTIVE = 1, an error is present, and a reaction that is determined by value VAL of DR\_REAC (FIG. 12) is implemented.

If DR\_REAC = DR\_OFF, execution jumps from S364 directly to S386, and no reaction occurs.

If DR\_REAC = DR\_n\_max, execution jumps from S366 to S368. In S368,

n\_const is set to 1 and "rotation speed calculation" function S626 (FIG. 6) is thereby informed that a constant rotation speed is now defined. In addition, rotation speed target value n\_s is set to a maximum rotation speed n\_max, pulse duty factor PWM\_TV of pulses 107A (FIG. 2) being set to 100%.

If DR\_REAC = DR\_n\_min in S370, then in S372 n\_const is set to 1 and rotation speed target value n\_s is set to a minimum rotation speed n\_min. This rotation speed is then specified to controller 70 as the rotation speed target value.

If DR\_REAC = DR\_n\_0 in S374, then in S376 n\_const is set to 1 and rotation speed target value n\_s is set to 0. For this purpose, transistors 114, 116, 118, 120 in FIG. 2 are made nonconductive so that motor 32 no longer receives current.

If DR\_REAC = DR\_n\_DIST in S378, then in S380 n\_const is set to 1 and rotation speed target value n\_s is set to error rotation speed n\_DIST which is defined by object n\_DIST (FIG. 8). This rotation speed n\_DIST can be conveyed to controller 70 as rotation speed target value n\_s.

If DR\_REAC = DR\_BRAKE in S382, then in S384 n\_const is set to 1, rotation speed target value n\_s is set to 0, and BRAKE is set to 1 in order to indicate to "RGL" function S610 (FIG. 6) that active braking is required. In this situation, for example, the two lower transistors 118, 120 of H-bridge 112 (FIG. 2) are made continuously conductive so that stator winding 128 is short-circuited, while upper transistors 114, 116 are opened.

After DR\_REAC has been checked in S364 through S382, S386 checks on the basis of DR\_AL (FIG. 12) whether an alarm is to be triggered in the event of an error.

If DR\_AL = DR\_AL\_ON, then in S388 state word AS\_DIST is set to 1, thereby informing alarm function S642 (FIG. 6) that an error is present and an alarm is to be triggered. The routine then branches to S390 (END).

COMMUNICATION FUNCTION (FIG. 24)

FIG. 24 shows a flow chart of "COMM" function S604 (FIG. 6). It controls input and output via bus 82. (Steps S402, S406, and S430 symbolize possible further program sections; i.e. in order to avoid unnecessary length, FIG. 24 usually represents only the portion of "COMM" function S604 that is important here.)

S404 is the beginning of portion PROCESS\_INSTR, in which instructions which the "COMM" function has received via a serial bus 82 (here an IIC bus), and which are located in INSTR, are executed.

The defined instructions begin with OC. OC\_GETDI means, for example, that state word DI\_STATE (FIG. 10) and error code DI\_CODE (FIG. 11) are polled from outside. S410 checks whether INSTR = OC\_GETDI. If Yes, the desired information is outputted onto IIC bus 82 using the instruction write\_IIC. OC\_DIDAT means that these are error data, and "2" means that two further arguments follow, namely DI\_STATE and DI\_CODE.

S414 is a comparison to determine whether INSTR = OC\_RES DI; if so, execution branches to S416. OC\_RES DI means that a reset of the error function must take place. Any errors saved in DI\_STATE, DI\_CODE must therefore be cleared. In S416, DI\_STATE and DI\_CODE are reset. DS\_CLASS and DS\_ACTIVE are reset simultaneously with DI\_STATE (FIG. 10).

S418 checks whether INSTR = OC\_GETOD; if so, execution branches to S420. OC\_GETOD means that the operating data of the "operating data function" (FIG. 15) are retrieved via bus 82. In S420, the instruction write\_IIC is used to write the operating data onto IIC bus 82. OC\_ODDAT means here that these are data of the "operating data function" (FIG. 15); and "4" means that four further arguments follow, namely OD\_UBM, OD\_TM, OD\_OH, and OD\_COM (cf. FIG. 15).

In S422, INSTR is compared to OC\_GETFIFO; if they are identical,



execution branches to S424. OC\_GETFIFO means that the next values are read out from the FIFO, which at each new error is filled with the error class, error code, and present operating data (FIG. 22). To achieve this, in S424 the next data are fetched from FIFO and written into variables TMP\_CLASS, TMP\_CODE, and TMP\_DATA. These data are then written onto bus 82 using write\_IIC. OC\_FIFO here indicates that these are data from the FIFO; and "3" indicates the number of additional parameters: TMP\_CLASS, TMP\_CODE, and TMP\_DATA.

S430 indicates possible further steps, and the COMM routine (which has a high priority of L3 as shown in FIG. 3) ends at S432.

In the present invention, therefore, a plurality of routines of different priorities (FIG. 6: L1 through L23) are provided, and if an error is identified as they are being executed, "emergency actions" are taken first. For example, if it is found at S226 in the A/D routine of FIG. 16 that an interruption exists in the line to sensor 152, then first of all, at S228, the corresponding error class and error code are saved, and then program NEW\_DIST of FIG. 22 is executed. If another error is not already present, DS\_ACTIVE is set therein. The error class and error code are saved in nonvolatile memory 74 in order to have permanent information about the nature of the error; similarly, all relevant operating data are saved in memory 74. The routine of FIG. 16 then initially continues to execute.

At some point the program in FIG. 6 then arrives at error function S640, which is depicted in FIG. 23.

The error function determines that DS\_ACTIVE is set; and in accordance with the parameters in word DIST\_REAC of object table 111 (these parameters are also located in RAM 97 while the motor is operating), a reaction to the error is then implemented. These parameters are depicted in FIG. 12. In the context of a fan, DR\_REAC (FIG. 12) will usually have the value 1, i.e. in the event of an error in temperature sensor 152,

the rotation speed of motor 32 is set to the maximum value in order to ensure reliable ventilation. This occurs in steps S366 and S368 of FIG. 23. Thus as soon as the connection to sensor 152 is interrupted, motor 32 is very quickly switched over to its maximum rotation speed.

If this is not desired, for example because the fan then becomes very loud, it is possible to save a specific rotation speed, e.g. 2500 rpm, in value n\_DIST of object table 111; the value selected for DR\_REAC in FIG. 12 is then "4", i.e. DR\_n\_DIST. In the event of an error the error function (FIG. 23) then goes to steps S378 and S380 and switches motor 32, when an error is detected, to a constant rotation speed of 2500 rpm. This rotation speed can be selected without restriction when the motor is parameterized, and can also be modified later if the user has suitable software.

In FIG. 12 it is similarly possible to define, by means of variable DR\_AL, whether or not an alarm is to be outputted. If DR\_AL has a value of 1 in this context, an alarm will be triggered by steps S386 and S388 of FIG. 23.

A motor can therefore easily be parameterized as to whether and how it will react if an error is identified. In any event, EEPROM 74 will have saved data containing the error class, error code, and relevant motor data at the moment of the error, for example operating hours, maximum operating voltage, maximum temperature, optionally the time of day and date, etc., thus making it much easier, or indeed possible at all, subsequently to analyze an error that has occurred.

The procedure upon occurrence of an error can be compared to the care given to an accident victim: First comes the paramedic, who applies a temporary bandage and writes down some brief information that is attached to the patient as a label. This corresponds to identification of the error and saving of the operating data, for example in the A/D routine (FIG. 16) or the error detection routine (FIG. 17). Then the paramedic sets flag DS\_ACTIVE = 1, and leaves the victim until an ambulance arrives.

The error function program (FIG. 23), i.e. the error monitoring routine, corresponds to the ambulance. It drives past at some time, recognizes the victim from flag DS\_ACTIVE = 1, picks him up, and treats him in accordance with the stored instructions.

A system of this kind is very open and can be expanded and modified in any direction, since modifications affect not the program but only the data in object table 111 (FIG. 8). In every situation, the paramedic arrives first and performs predetermined actions; then the ambulance arrives and also performs predetermined actions. And those predetermined actions can be parameterized in memory 74.

Another advantage of the invention is that on the basis of objective data, a decision can be made as to whether a motor has reached the end of its service life and should be replaced as a precaution.

A suitable time interval for transferring the operating data variables into the nonvolatile memory is, for example, a fraction of an hour, e.g. 10 minutes, 20 minutes, 30 minutes, or the like. The transfer can thus take place relatively infrequently and therefore does not demand a great deal of calculation time. The accuracy of the data saved in the nonvolatile memory is nevertheless good, since electric motors usually run uninterruptedly for a long period, during which their operating data change only slightly.

A further possibility for saving the present variables as old operating data values in nonvolatile memory 74 is available in the context of microprocessors 23 which can still execute shutdown routines during a reset, occurring, for example, in the event of a power failure or when electric motor 32 is switched off. The present variables can be saved to the nonvolatile memory in these shutdown routines, and more-frequent intermediate saving is no longer necessary.

Many modifications and variants are, of course, possible within the scope of the present invention.